

libEnsemble: A Library for Managing Ensembles of Calculations

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▶ libEnsemble uses a manager to allocate work to various workers.

▶ A libEnsemble worker is the smallest indivisible unit to perform



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sim_f: Evaluates a simulation (i.e., user-defined function) using input
 defined by gen_f.

alloc_f: Decides whether (or not) sim_f or gen_f should be called (and with what input/resources) as workers become available.



libEnsemble dependencies

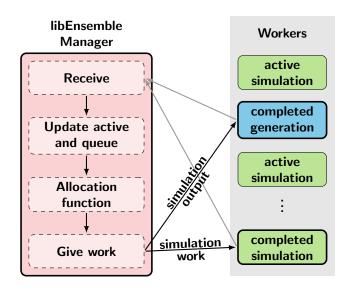
▶ Python 2.7, 3.4 or above.

► An MPI 1.x/2.x/3.x implementation (e.g., MPICH or OpenMPI) built with shared/dynamic libraries.

mpi4py v2.0.0 or above.

▶ NumPy v1.10 or above.





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- (*) Execution on multiple LCFs
- (*) Thousands of concurrent workers

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- ▶ Don't have to write your own kills, just complete libEnsemble templates.
- ► Want to add concurrency to a generator (e.g., multiple local optimizers.)

Why libEnsemble and not...?

Dakota: Few simulation optimization generators.

► C++

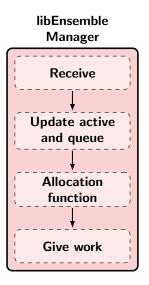
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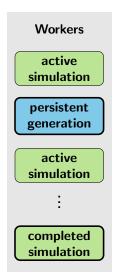
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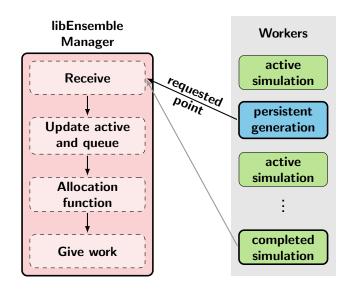
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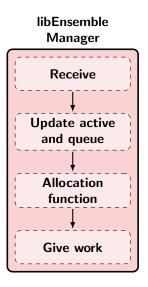
Swift: (the parallel scripting language)

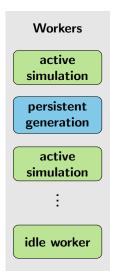
- "Can run million programs, thousands at a time, launching hundreds per second."
- Require writing your generators in Swift's scripting language.
- Difficult to tightly couple generation of inputs and future/active running simulations.

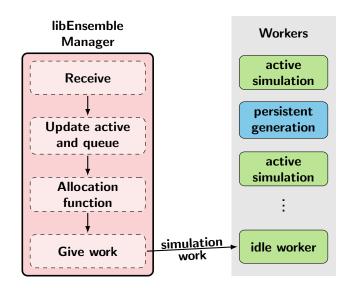


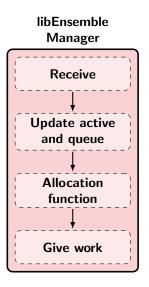


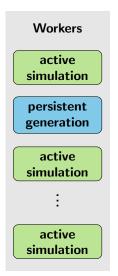


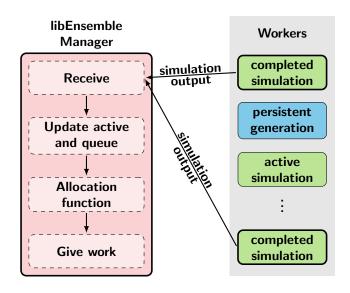


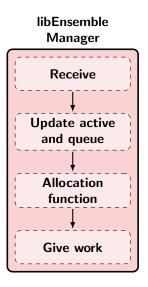


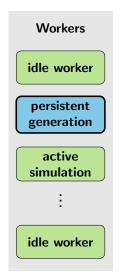


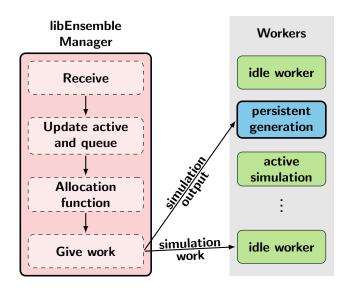












Timing libEnsemble overhead

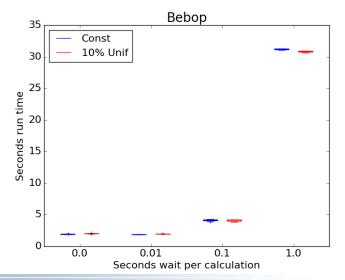
Time for libEnsemble to sample/evaluation 30×(workers) points

Tons of system-dependent caveats

Timing libEnsemble overhead

Time for libEnsemble to sample/evaluation $30 \times$ (workers) points

32 nodes \times 36 cores = 1152-1 workers

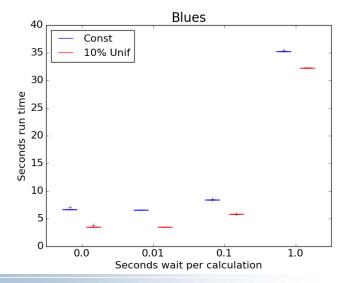




Timing libEnsemble overhead

Time for libEnsemble to sample/evaluation $30 \times (workers)$ points

64 nodes \times 16 cores = 1024-1 workers





- ▶ A user wants to optimize a function that depends on a simulation.
- ► The simulation is already using parallel resources, but not a large fraction of some computer.
- ▶ libEnsemble can coordinate the concurrent evaluation of the simulation sim_f at various parameter values and gen_f would return candidate parameter values (possibly after each sim_f output).

- ► A user has a gen_f that produces different meshes to be used within a sim_f.
- ► Given the sim_f output, gen_f will refine a mesh or produce a new mesh.
- ▶ libEnsemble can ensure that the calculated meshes can be used by multiple simulations without requiring movement of data.

- ► A user wants to evaluate a simulation sim_f at parameters sampled from a set of parameter values.
- Many parameter sets will cause the simulation to fail.
- libEnsemble can stop unresponsive evaluations, and recover computational resources for future evaluations.
- gen_f can update the sampling after discovering regions where evaluations of sim_f fail.

- ► A user has a simulation sim_f that requires calculating multiple expensive quantities, some of which depend on other quantities.
- ▶ libEnsemble can observe intermediate quantities in order to stop related calculations and preempt future calculations associated with a poor parameter values.

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Naturally, combinations of use cases is supported as well.



Problem setup

▶ We want to identify distinct, "high-quality", local minimizers of

minimize
$$f(x)$$

 $1 \le x \le u$
 $x \in \mathbb{R}^n$

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- Derivatives of f may or may not be available.
- ► The simulation *f* is likely using parallel resources, but it does not utilize the entire machine.

Global optimization is difficult

Theorem (Törn and Žilinskas, Global Optimization, 1989)

An algorithm converges to the global minimum of any continuous f on a domain \mathcal{D} if and only if the algorithm generates iterates that are dense in \mathcal{D} .



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The theory can be more than merely checking that a method generates iterates which are dense in the domain.



Given some local optimization routine \mathcal{L} :

Algorithm 1: General Multistart

for k = 1, 2, ... do

Evaluate f at N points drawn from \mathcal{D}

Start ${\mathcal L}$ at some set (possibly empty) of previously evaluated points



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- ▶ If resources are limited, how should points from each run receive priority?

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- Which points should start runs?
- If resources are limited, how should points from each run receive priority?
- ▶ Ideally, only one run is started for each minima.
- **Exploring by sampling.** Refining with \mathcal{L} .

Multi-Level Single Linkage

- ▶ $f \in C^2$, with local minima in the interior of \mathcal{D} , and the distance between these minima is bounded away from zero.
- $ightharpoonup \mathcal{L}$ is strictly descent and converges to a minimum (not a stationary point).

$$r_k = \frac{1}{\sqrt{\pi}} \sqrt[n]{\Gamma\left(1 + \frac{n}{2}\right) \operatorname{vol}(\mathcal{D})} \frac{\sigma \log kN}{kN}$$
 (1)

Theorem

If $r_k \to 0$, all local minima will be found almost surely.

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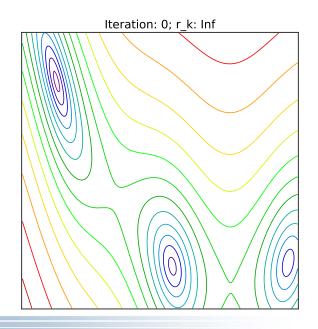
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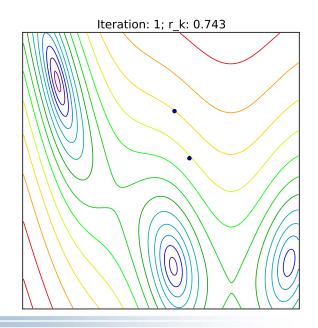
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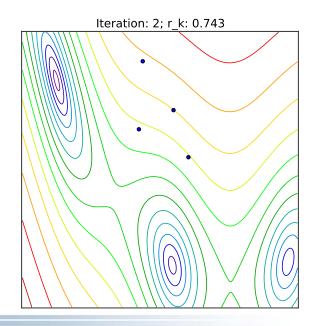
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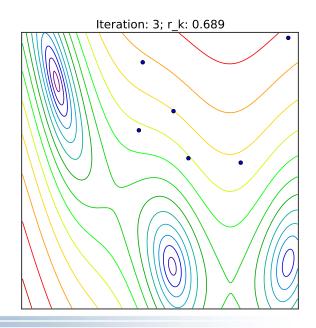
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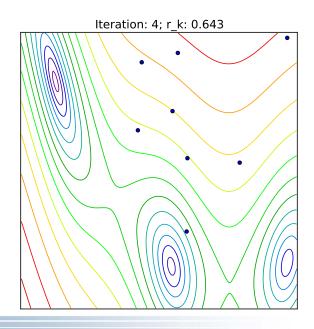
If r_k is defined by (1) with $\sigma > 4$, even if the sampling continues forever, the total number of local searches started is finite almost surely.

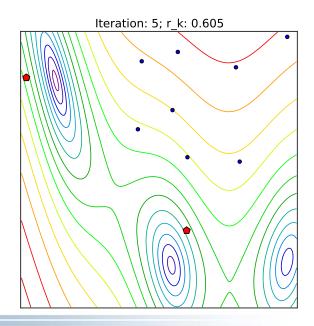


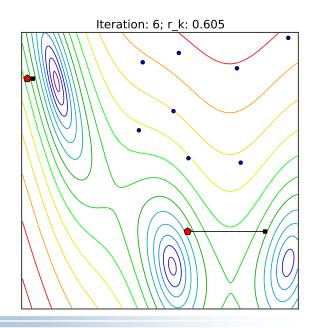


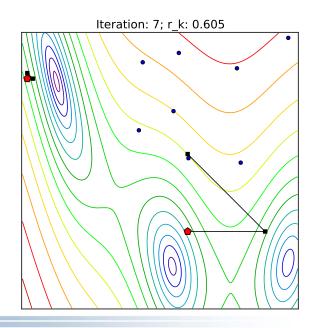


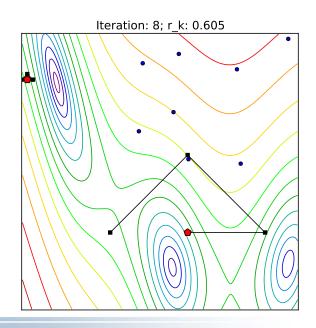


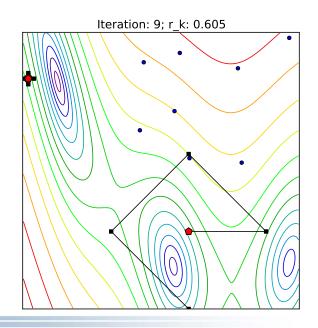


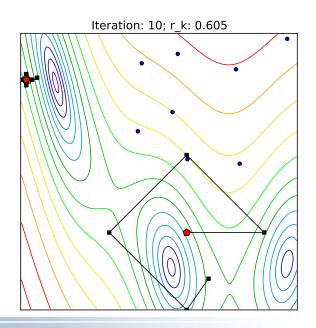


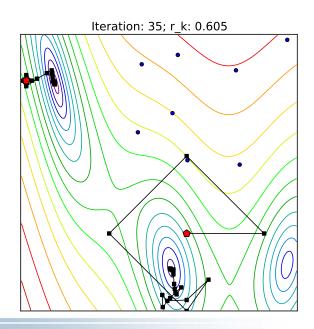


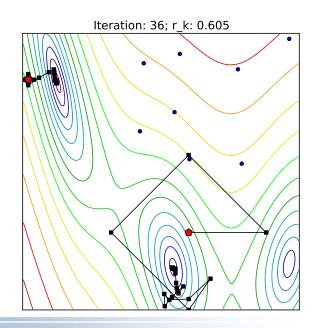


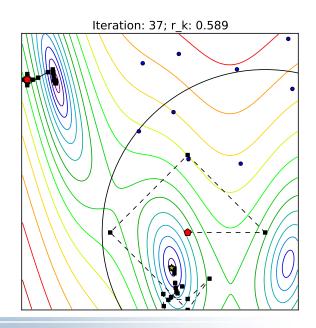


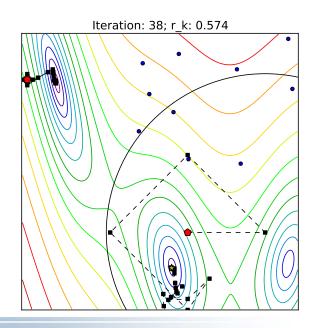


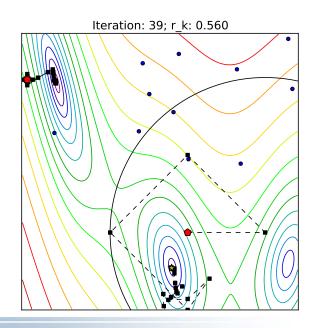


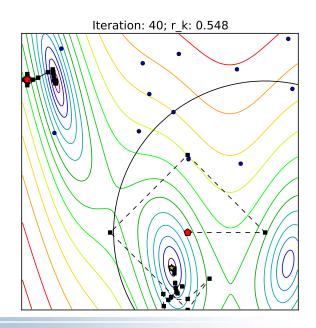


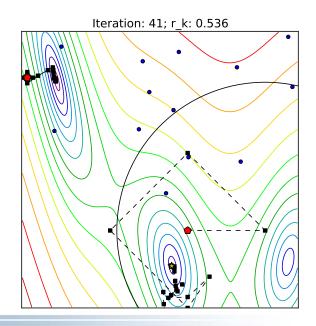


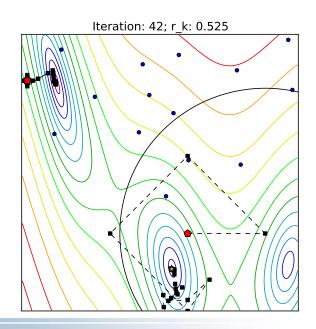


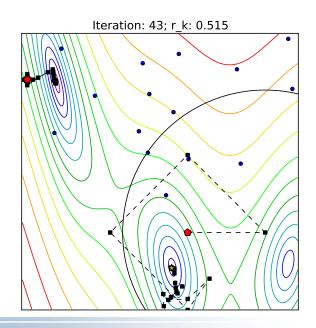


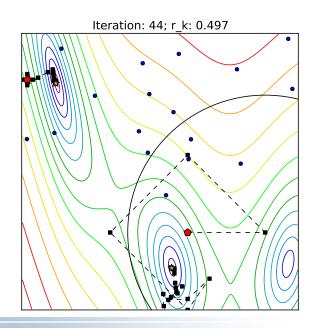


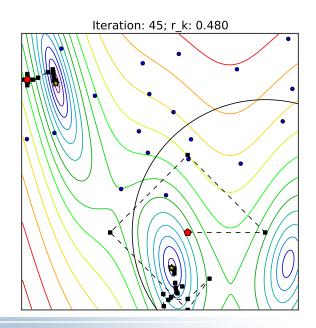


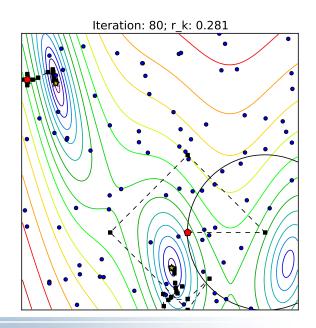


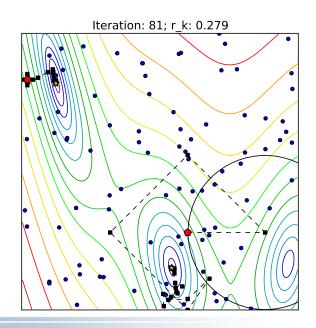


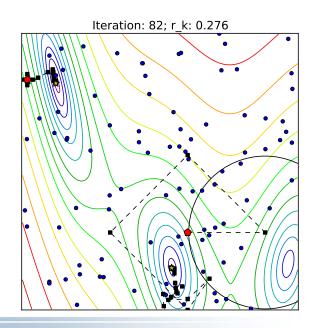


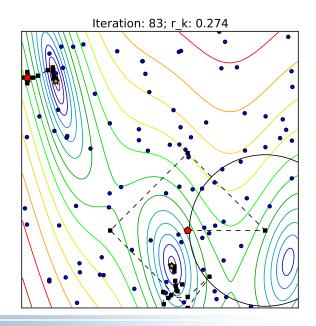


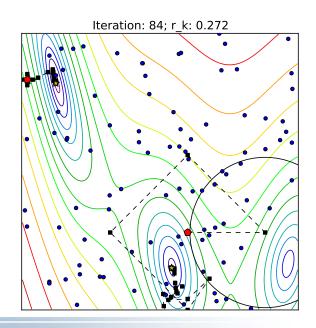


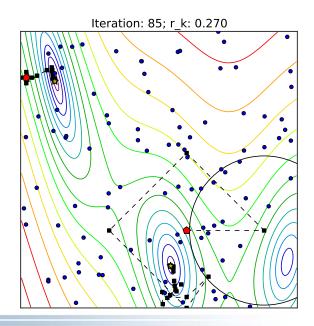


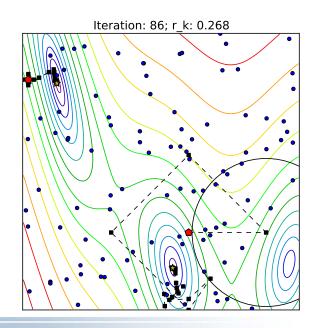


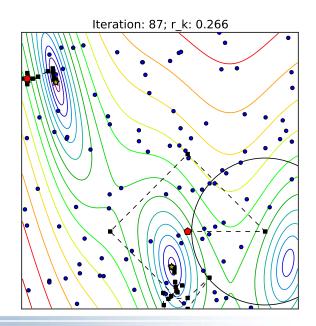


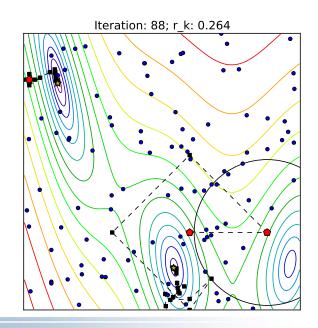


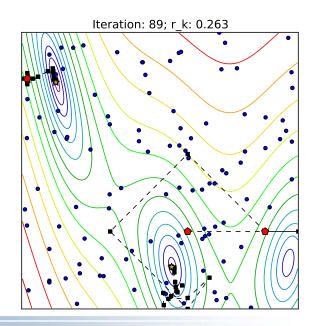


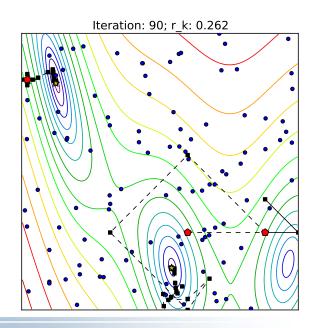


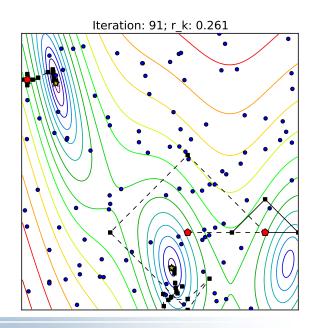


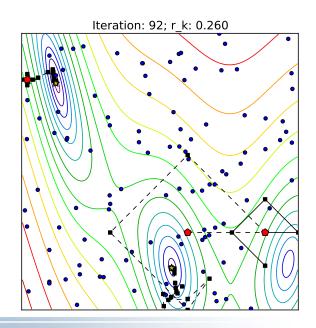


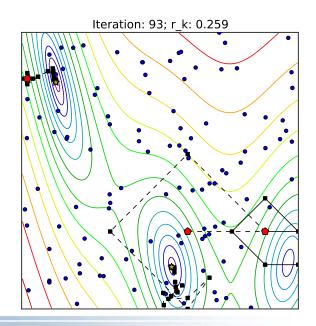


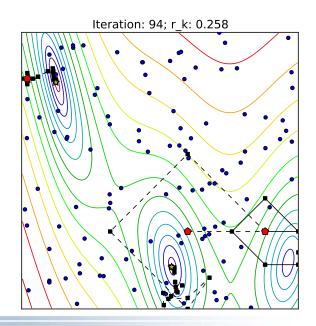


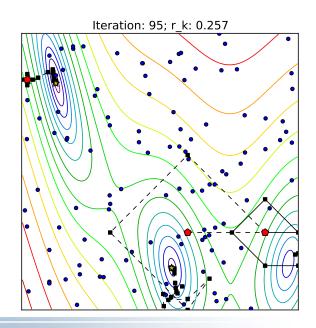


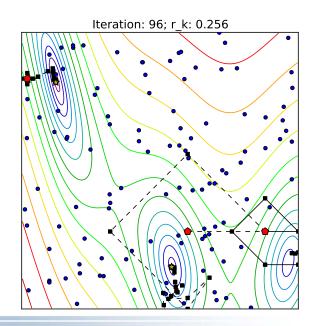


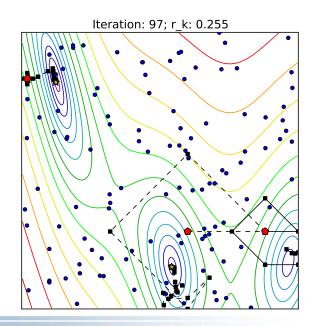


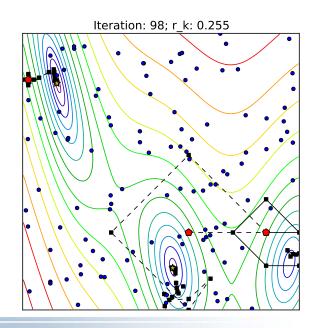


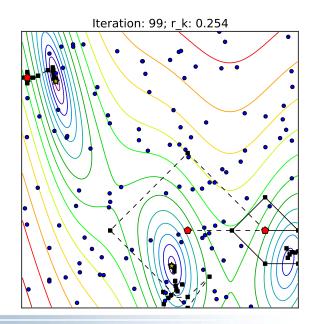












Properties of the local optimization method

Necessary:

- ► Honors a starting point
- ► Honors bound constraints



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Possibly beneficial:

- Can return multiple points of interest
- Reports solution quality/confidence at every iteration
- Can avoid certain regions in the domain
- Uses a history of past evaluations of f
- ▶ Uses additional points mid-run

Measuring Performance

Let X^* be the set of all local minima of f.

Let $f_{(i)}^*$ be the *i*th smallest value $\{f(x^*)|x^* \in X^*\}$. Let $x_{(i)}^*$ be the element of X^* corresponding to the value $f_{(i)}^*$.

The global minimum has been found at a level $\tau > 0$ at batch k if an algorithm it has found a point \hat{x} satisfying:

$$f(\hat{x}) - f_{(1)}^* \le (1 - \tau) \left(f(x_0) - f_{(1)}^* \right),$$

where x_0 is the starting point for problem p.



Measuring Performance

Let X^* be the set of all local minima of f.

Let $f_{(i)}^*$ be the *i*th smallest value $\{f(x^*)|x^* \in X^*\}$. Let $x_{(i)}^*$ be the element of X^* corresponding to the value $f_{(i)}^*$.

The j best local minima have been found at a level $\tau > 0$ at batch k if:

$$\left|\left\{x_{(1)}^*, \dots, x_{(\underline{j}-1)}^*\right\} \bigcap \left\{x_{(i)}^* : \exists x \in \mathcal{H}_k \text{ with } \left\|x - x_{(i)}^*\right\| \leq r_n(\tau)\right\}\right| = \underline{j} - 1$$

$$\left|\left\{x_{(\underline{j})}^*, \dots, x_{(\overline{j})}^*\right\} \bigcap \left\{x_{(i)}^* : \exists x \in \mathcal{H}_k \text{ with } \left\|x - x_{(i)}^*\right\| \leq r_n(\tau)\right\}\right| \geq \underline{j} - \underline{j} + 1,$$

where j and \bar{j} are the smallest and largest integers such that

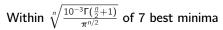
$$f_{(j)}^* = f_{(j)}^* = f_{(j)}^*$$
 and where $r_n(\tau) = \sqrt[n]{rac{ au \operatorname{vol}(\mathcal{D})\Gamma(rac{n}{2}+1)}{\pi^{n/2}}}$.

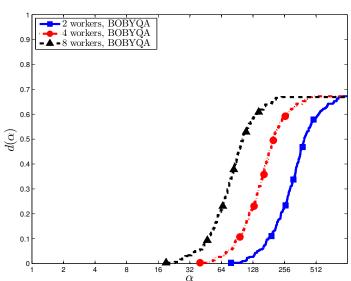


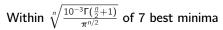
Problems considered

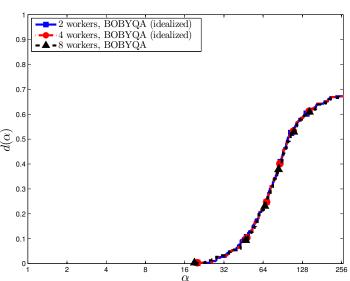
GKLS problem generator [Gaviano et al., "Algorithm 829" (TOMS, 2003)]

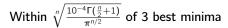
- ▶ 600 synthetic problems with known local minima
- ▶ n = 2, ..., 7
- ▶ 10 local minima in the unit cube with a unique global minimum
- ▶ 100 problems for each dimension
- 5 replications (different seeds) for each problem
- ▶ 5000 evaluations

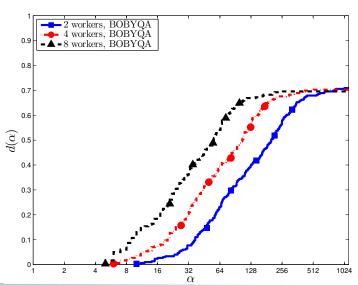


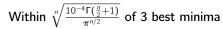


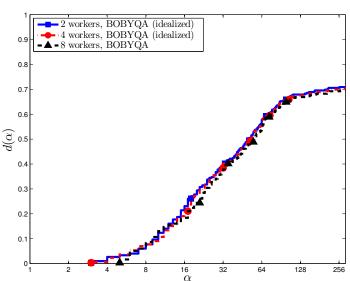




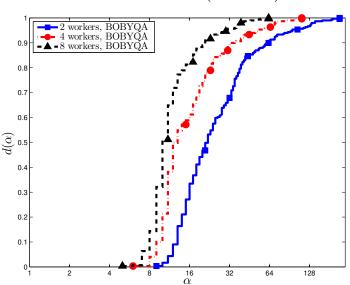




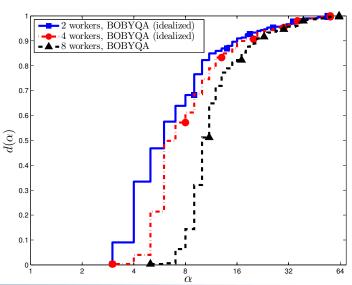




$$f(x) - f_{(1)}^* \le (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$



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Closing Remarks

libEnsemble is a software library to coordinate the concurrent evaluation of calculations.

We have a growing set of use cases and examples.

▶ Let us know if you have examples you'd like to see.

